# CS275 Discrete Mathematics 

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## Goal for labs

- Review contents
- Practice for homeworks/tests
- Answer questions
- Help you better understand the course \& get the grade you aimed


## Logic and Proof Section 1.1-1.6

## What is a proposition?

- A proposition is a declarative statement that is True or False but not both.
- E.g., Tony is original from China.

Negation of a proposition ( $\neg \mathrm{p}$ )

- $\quad \mathrm{p}$ : it is not the case that p
- E.g.,
- $p=$ true, $\neg p=$ false
- $\quad \mathrm{P}=$ "Today is Wed.", $\neg \mathrm{p}=$ "Today is NOT

Wed."

## Conjunction of $p$ and $q(p \wedge q)$

- The conjunction $p \wedge q$ ( $p$ and $q$ ) is true if both $p$ and $q$ are true; otherwise it is false.
- E.g.,

$$
\begin{aligned}
& \text { - } \mathrm{p}=\text { "Today is Wed.", } \mathrm{q}=\text { "Today is 01/01." } \\
& \text { - } \mathrm{p} \wedge \mathrm{q}=\text { ? }
\end{aligned}
$$

## Disjunction of $p$ and $q(p \vee q)$

- The disjunction $\mathrm{p} \vee \mathrm{q}(\mathrm{p}$ or q$)$ is false if both p and q are false; otherwise, it is true.
- E.g.,
- $p=$ "Today is Wed.", $q=$ "Today is 01/01."
- $\mathrm{p} \vee \mathrm{q}=$ ?


## Conditional statement ( $p \rightarrow q$ )

- The conditional statement $\mathrm{p} \rightarrow \mathrm{q}$ (if p then q ) is false when p is true and q is false; otherwise, it is true.
- E.g., $\mathrm{p}=$ "If I have a keyboard", $\mathrm{q}=$ "I can type"
- If $p$ is true, $q$ is true, $p \rightarrow q$ is true
- If $p$ is false, $q$ is false
- "If I DON'T have a keyboard, I CAN'T type"
- Could be! Thus, $\mathrm{p} \rightarrow \mathrm{q}=$ true
- If $p$ is false, $q$ is true,
- "If I don't have a keyboard, I still type", thus, $p \rightarrow q$ is true
- If $p$ is true, $q$ is false,
- "If I have a keyboard, I cannot type"
- Why?! Thus, $p \rightarrow q=$ false


## Conditional statement $(p \rightarrow q)$

- $\mathbf{p} \rightarrow \mathbf{q} \equiv \neg \mathbf{p} \vee \mathrm{q}$
- E.g., p = "If I have a keyboard", q = "I can type"
- If $p$ is true, $q$ is true
- $\quad$ p V q: "if I don't have a keyboard, I can type" TRUE
- If $p$ is false, $q$ is false
- $\quad$ p $\vee$ q: "if I have have a keyboard, I cannot type" FALSE
- If $p$ is false, $q$ is true
- $\quad \neg \mathrm{p} \vee \mathrm{q}$ : "if I don't have a keyboard, I still can type" TRUE
- If $p$ is true, $q$ is false
- $\quad$ p V q: "if I don't have a keyboard, I cannot type" TRUE


## Prove $p \rightarrow q \equiv \neg p \vee q$ (using truth table)

| $\mathbf{p}$ | $\mathbf{q}$ | $\mathbf{p} \rightarrow \mathbf{q}$ | $\boldsymbol{\sim p}$ | $\boldsymbol{\sim} \mathbf{p} \mathbf{q}$ |
| :---: | :---: | :---: | :---: | :---: |
| T | T | T | F | T |
| T | F | F | F | F |
| F | T | T | T | T |
| F | F | T | T | T |

## Converse of $\mathbf{p} \rightarrow \mathbf{q}$ \& Contraposition of $\mathbf{p} \rightarrow \mathbf{q}$

- Converse of $p \rightarrow q$ is $q \rightarrow p$
- The contraposition of $p \rightarrow q$ is $\neg q \rightarrow \neg p$
- $p \rightarrow q \equiv \neg q \rightarrow \neg p$

| $\mathbf{p}$ | $\mathbf{q}$ | $\mathbf{p} \rightarrow \mathbf{q}$ | $\boldsymbol{\neg q}$ | $\boldsymbol{\neg p}$ | $\boldsymbol{\neg q} \rightarrow \boldsymbol{\mathbf { p }}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| T | T | T | F | F | T |
| T | F | F | T | F | F |
| F | T | T | F | T | T |
| F | F | T | T | T | T |

## Biconditional statement ( $\mathbf{p} \leftrightarrow \mathbf{q}$ )

- A biconditional statement $p \leftrightarrow q$ (if $p$ and only if $q$ ) is true if both operands are true or both operands are false

| $\boldsymbol{p}$ | $\boldsymbol{q}$ | $\boldsymbol{p} \leftrightarrow \boldsymbol{q}$ |
| :---: | :---: | :---: |
| T | T | T |
| T | F | F |
| F | T | F |
| F | F | T |

Exercise: Prove $p \leftrightarrow q \equiv(p \wedge q) \vee(\neg p \wedge \neg q)$ using truth table

| $\mathbf{p}$ | $\mathbf{q}$ | $\mathbf{p} \leftrightarrow \mathbf{q}$ | $\mathbf{p} \wedge \mathbf{q}$ | $\boldsymbol{\imath p}$ | $\boldsymbol{q} \mathbf{q}$ | $\neg \mathbf{p} \wedge \mathbf{q}$ | $(\mathbf{p} \wedge \mathbf{q}) \vee(\neg \mathbf{p} \wedge \neg \mathbf{q})$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| T | T | T | T | F | F | F | T |
| T | F | F | F | F | T | F | F |
| F | T | F | F | T | F | F | F |
| F | F | T | F | T | T | T | T |

